

# PHPE 308M/PHIL 209F

## Fairness

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September 10, 2025

# Preferences in the Ultimatum Game

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- ▶ Players are assumed to have utility functions (a function that maps outcomes to real numbers) *representing* their preferences over the outcomes:  
The utilities for the outcome  $(x_P, x_R)$  are  $u_P(x_P, x_R)$  and  $u_R(x_P, x_R)$ .

## Social Preferences

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Social preferences are examples of **outcome-based preferences**: utility functions that depend only on:

1. the individuals involved in the interaction and
2. the monetary payoffs associated with each strategy profile.

## Inequality Aversion: Fehr and Schmidt Utility Function

$$u_P(x_P, x_R) = x_P - \alpha_P \max(x_R - x_P, 0) - \beta_P \max(x_P - x_R, 0)$$

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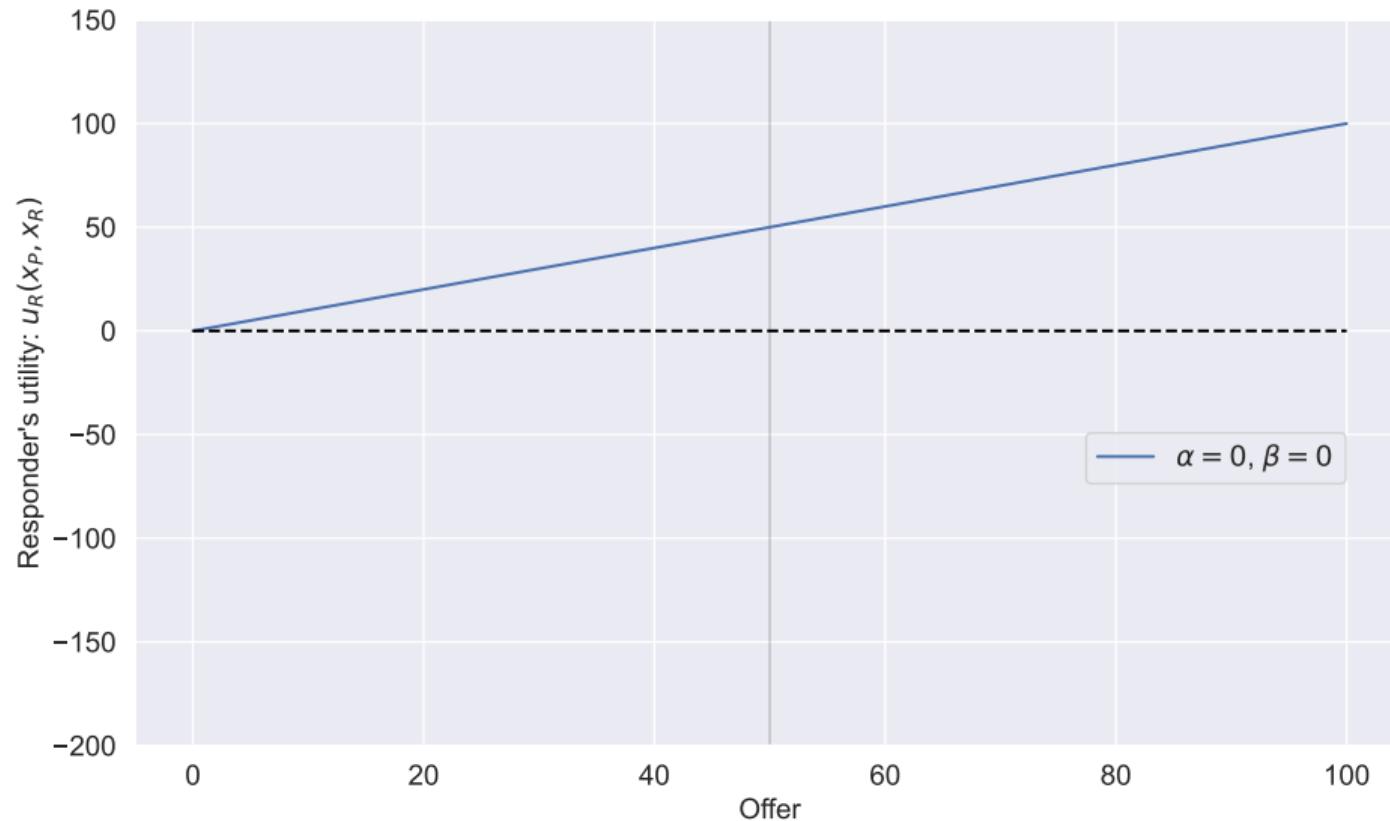
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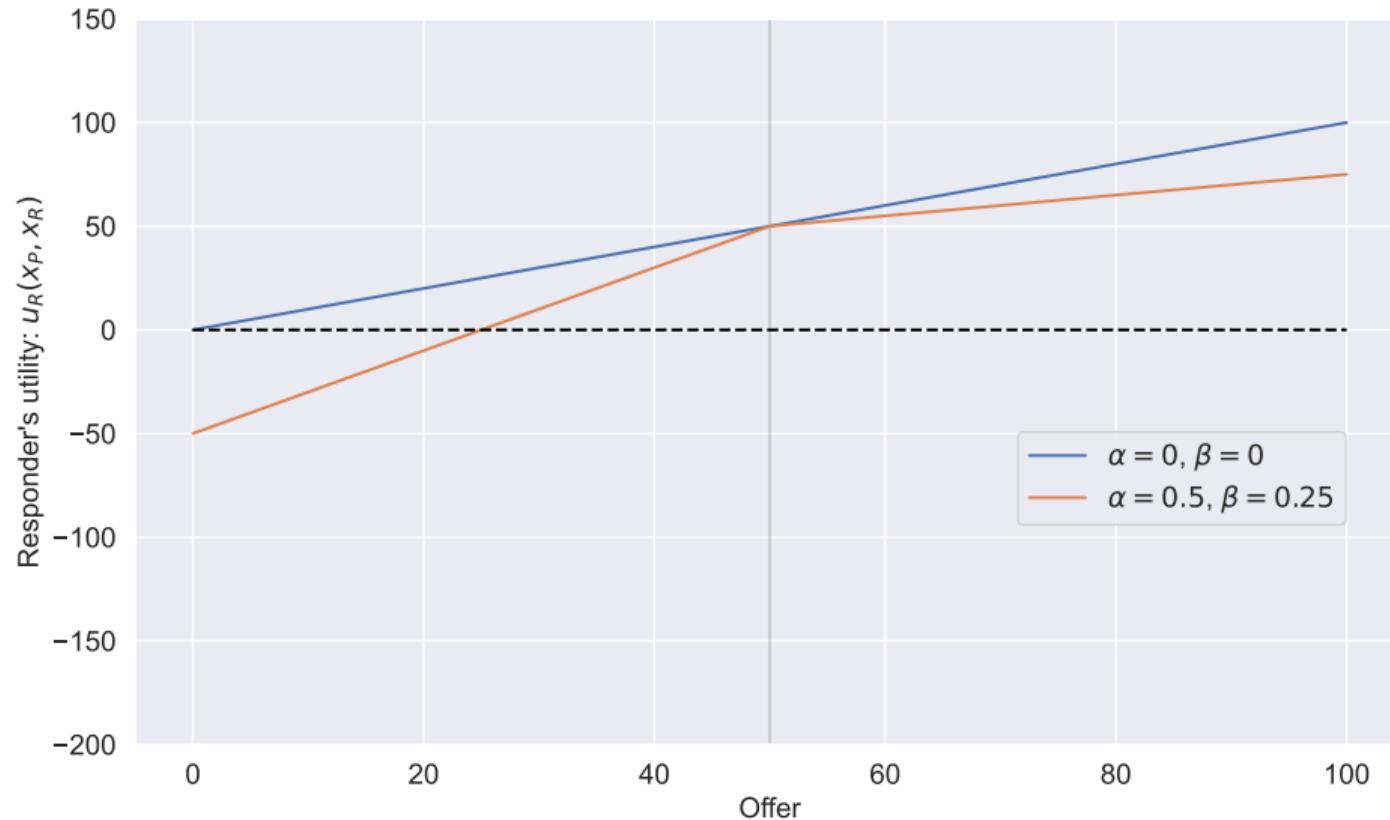
- ▶  $\alpha_i$  is  $i$ 's 'envy' weight and  $\beta_i$  is  $i$ 's 'guilt' weight
- ▶  $0 < \beta_i < \alpha_i$ : indicates that people dislike inequality against them more than they do inequality favoring them.
- ▶  $\beta_i < 1$ : agents do not suffer terrible guilt when she is in a relatively good position. For example, a player would prefer getting more without affecting other people's payoff even though that increases inequality.

Ernst Fehr and Klaus M. Schmidt (1999). *A theory of fairness, competition, and cooperation*. The Quarterly Journal of Economics, 114(3), pp. 817 - 868.

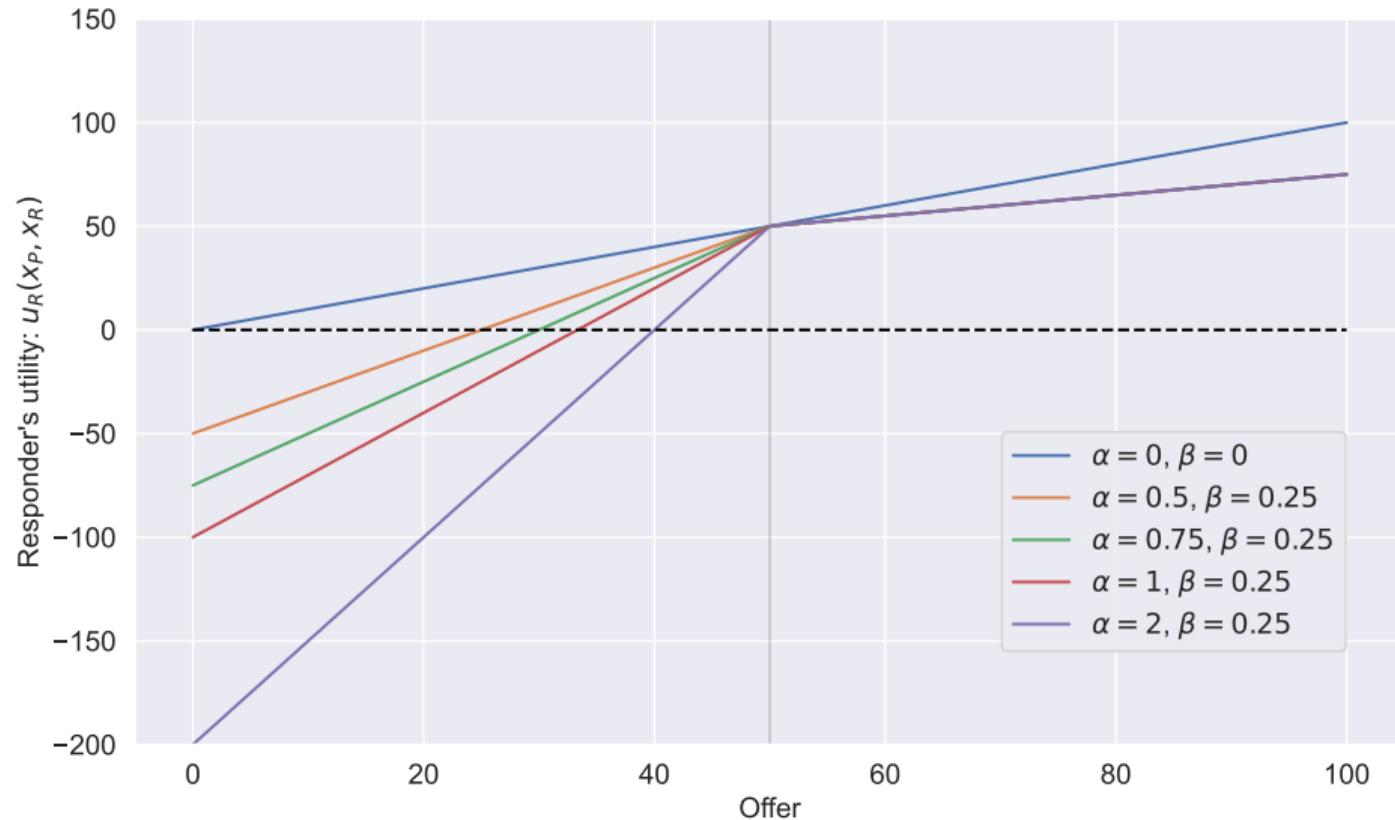
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## The Responder's Utility

Suppose that the total amount to be distributed is  $M$  and  $y$  is how much  $P$  offers to  $R$ . So,  $x_P = M - y$  and  $x_R = y$ .

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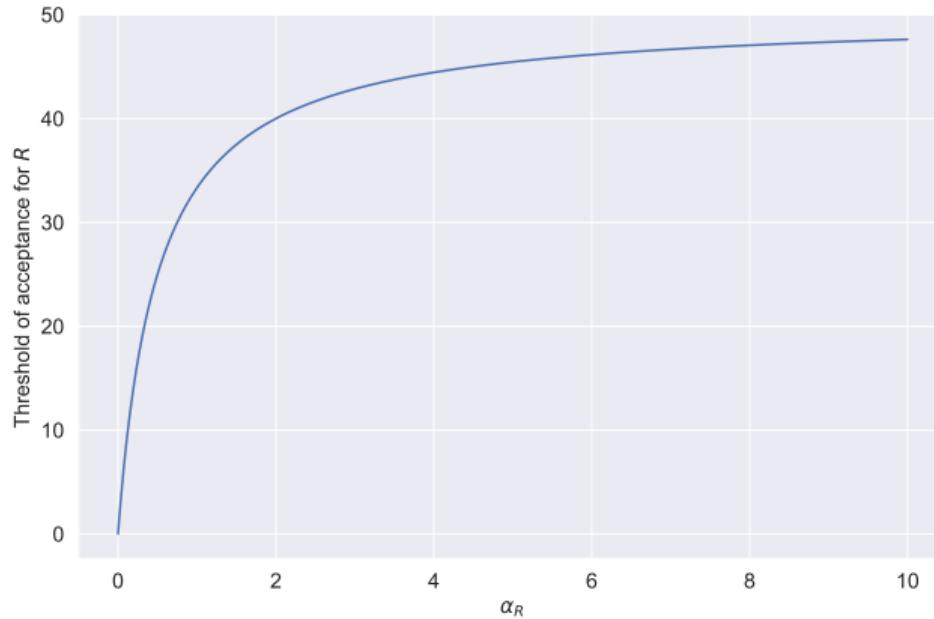
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Then,  $R$  should accept provided that  $u_R(x_P, x_R) > 0$ . Solving, for  $y$ , we get:

$$y > \frac{\alpha_R M}{1 + 2\alpha_R}$$

## Acceptance Threshold for $R$



If  $\alpha_R$  is close to zero — which indicates that  $R$  does not care much about being treated unfairly — the responder will accept very low offers. On the other hand, if  $\alpha_R$  is sufficiently large, the offer has to be close to a half to be accepted.

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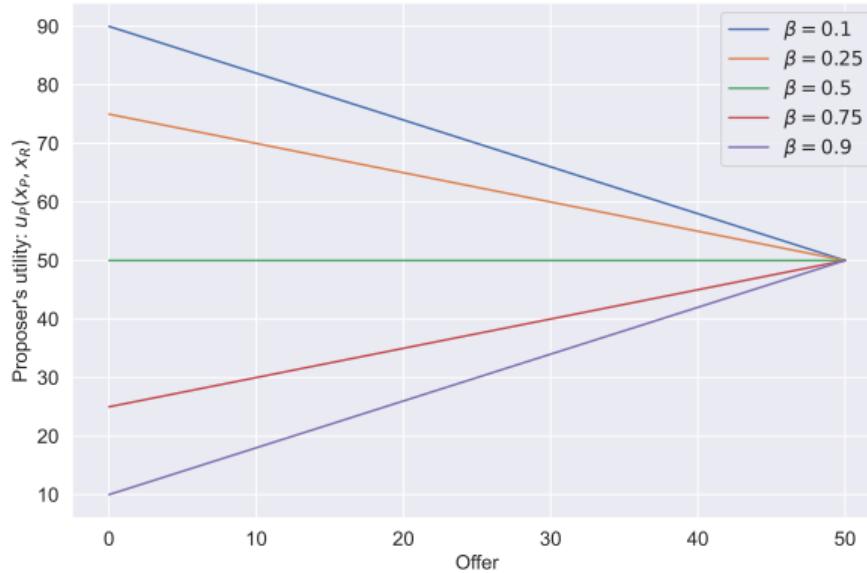
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# Proposer's Best Offer



If  $\beta_P > 1/2$  — if the proposer feels sufficiently guilty about treating others unfairly — then the best choice is to offer  $M/2$ . If  $\beta < 1/2$ , then the best offer is the minimum one that would be accepted, (a little bit more than  $\alpha_R M / (1 + 2\alpha_R)$ ). If  $\beta = 1/2$ , then any (acceptable) offer is best.

## The Fehr-Schmidt Utility Function

As noted by Fehr and Schmidt, the model allows for the fact that individuals are heterogeneous. Different  $\alpha$ s and  $\beta$ s correspond to different types of people. Although the utility functions are common knowledge, the exact values of the parameters are not. The proposer, in most cases, is not sure what type of responder she is facing.

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The experimental data suggest that for many proposers, either  $\beta_P$  is large ( $\beta_P > 1/2$ ) or they estimate the responder's  $\alpha_R$  to be large.

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One of the major weaknesses of this model, however, is that it has a **consequentialist bias**: players only care about final distributions of outcomes, not about how such distributions come about.

Shortly after the explosion of inequity aversion models, several economists observed that some decision-makers appear to act in a way that *increases* inequity, if this increase results in an increase in the total payoff of the participants.... This observation is hard to reconcile with inequity-aversion models, and suggests that people not only prefer to minimize inequity, but also prefer to maximize social welfare.

James Andreoni and John Miller (2002). *Giving According to GARP: An Experimental Test of the Consistency of Preferences for Altruism*. *Econometrica*, 70(2), pp. 737-753.

A. Falk, E. Fehr, and U. Fishbacher (2003). *On the nature of fair behavior*. Economic Inquiry, 41(1), pp. 20 - 26.

A. Festré (2019). *On the Nature of Fair Behaviour: Further Evidence*. Homo Oeconomicus, 36, pp. 193 - 207.

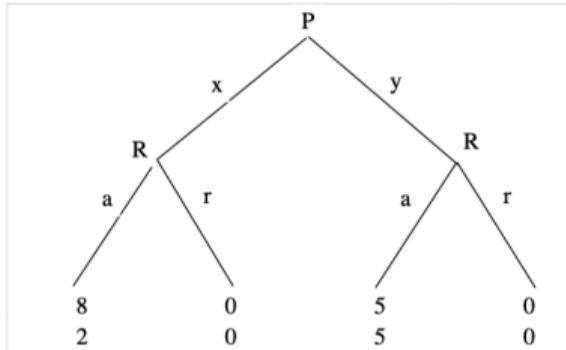
...identical offers in an ultimatum game trigger vastly different rejection rates depending on the other offers available to the proposer. In particular, a given offer with an unequal distribution of material payoffs is much more likely to be rejected if the proposer could have proposed a more equitable offer than if the proposer could have proposed only more unequal offers.

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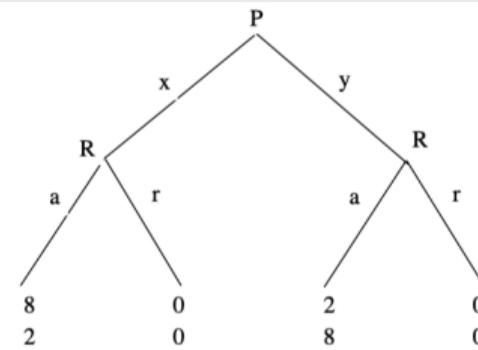
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Each of 90 experimental subjects participated in four different games. In all games the proposer P is asked to divide 10 points between himself and the responder R, who can either accept or reject the offer. Accepting the offer leads to a payoff distribution according to the proposer's offer. A rejection implies zero payoffs for both players.

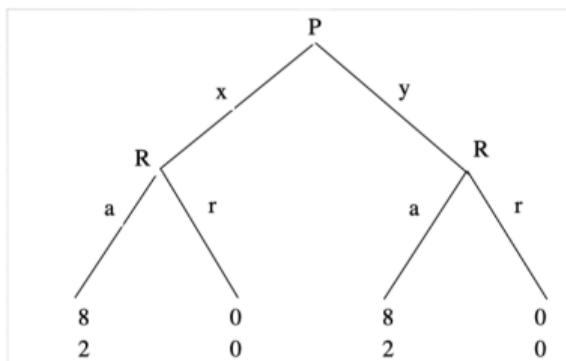
**Figure 1: The mini ultimatum games**



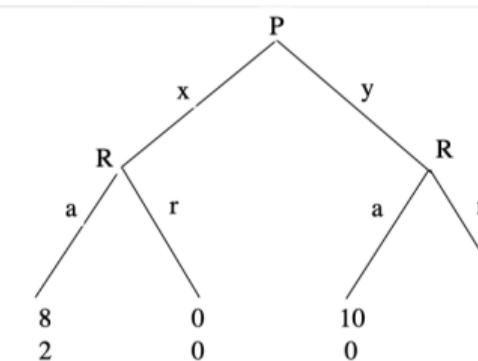
(a) (5/5)-game



(b) (2/8)-game



(c) (8/2)-game



(d) (10/0)-game

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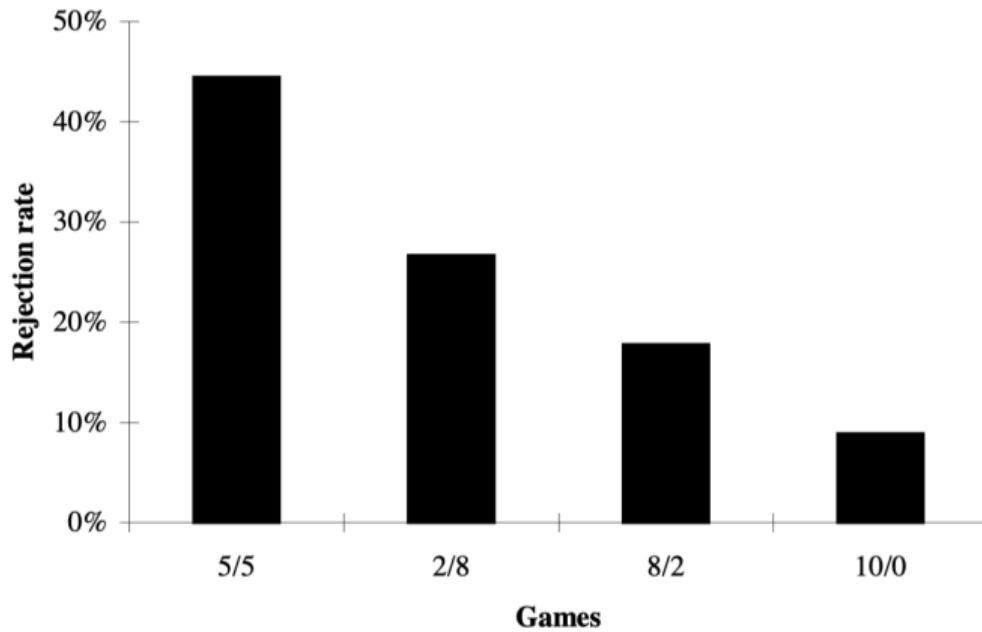
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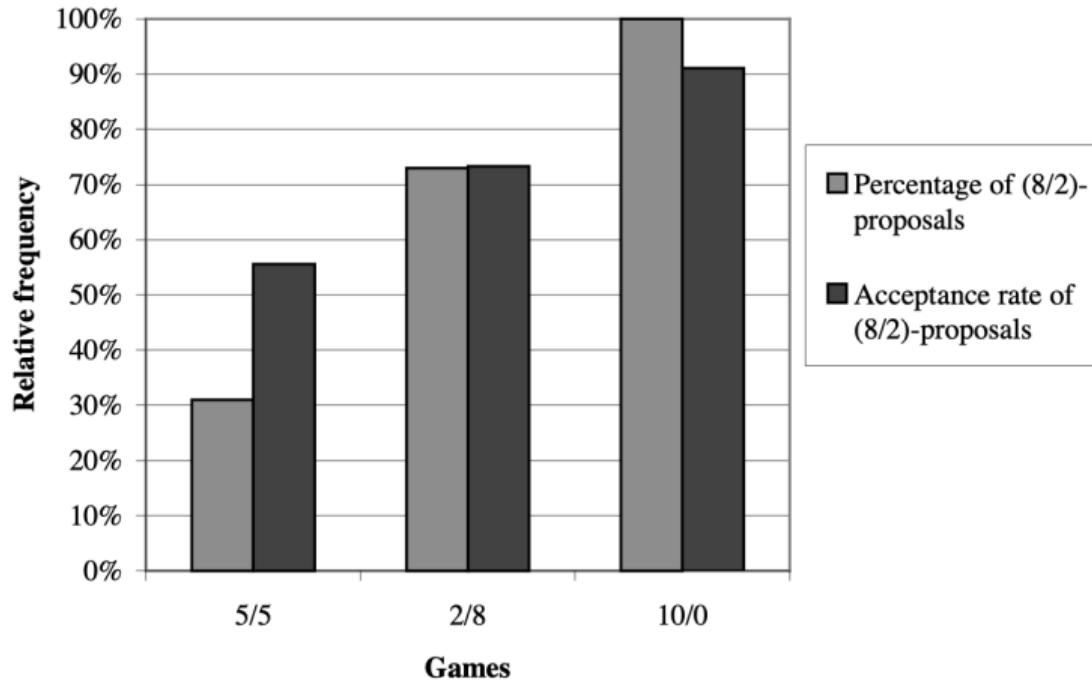
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After the end of the fourth game subjects received a show-up fee of plus their earnings from the experiment (about \$23 was at stake).

**Figure 2**  
**Rejection rate of the (8/2)-offer across games**



**Figure 3:**  
**Percentage and acceptance of (8/2) proposals**



The results of our experiment clearly show that the same action by the proposer in a miniultimatum game triggers very different responses depending on the alternative action available to the proposer. **This suggests that responders do not only take into account the distributive consequences of the action by the proposer but also the intention that is signaled by the action.**

- ✓ Social Preferences: Outcome-based Preferences
- ▶ Moral Preferences
- ▶ Reasoning based on Norms

J. Halpern, V. Capraro and M. Perc (2022). *From outcome-based to language-based preferences*. Journal of Economic Literature.

# Nash Bargaining Game

Brian Skyrms (2012). Chapters 1 & 2 in *Evolution of the Social Contract*. Cambridge University Press.

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1. 0: Demand nothing
2. 1/3: Demand 1/3 of \$1 (so 33 cents)
3. 1/2: Demand 1/2 of \$1 (so 50 cents)
4. 2/3: Demand 2/3 of \$1 (so 66 cents)
5. 1: Demand \$1

## Simplified Nash Bargaining Game

		Player 2				
		0	1/3	1/2	2/3	1
Player 1	0	0, 0	$0, \frac{1}{3}$	$0, \frac{1}{2}$	$0, \frac{2}{3}$	0, 1
	$\frac{1}{3}$	$\frac{1}{3}, 0$	$\frac{1}{3}, \frac{1}{3}$	$\frac{1}{3}, \frac{1}{2}$	$\frac{1}{3}, \frac{2}{3}$	0, 0
	$\frac{1}{2}$	$\frac{1}{2}, 0$	$\frac{1}{2}, \frac{1}{3}$	$\frac{1}{2}, \frac{1}{2}$	0, 0	0, 0
	$\frac{2}{3}$	$\frac{2}{3}, 0$	$\frac{2}{3}, \frac{1}{3}$	0, 0	0, 0	0, 0
	1	1, 0	0, 0	0, 0	0, 0	0, 0

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	$\frac{1}{2}$	$\frac{1}{2}, 0$	$\frac{1}{2}, \frac{1}{3}$	$\frac{1}{2}, \frac{1}{2}$	0, 0	0, 0
	$\frac{2}{3}$	$\frac{2}{3}, 0$	$\frac{2}{3}, \frac{1}{3}$	0, 0	0, 0	0, 0
	1	1, 0	0, 0	0, 0	0, 0	0, 0